

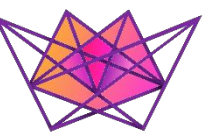
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Bra-ket notation (Dirac, 1939)

- V – vector (Hilbert) space, \mathbb{F} – field
- $|\psi\rangle$ – *pure state* (vector, or operator $\mathbb{F} \rightarrow V$) – ket
- $\langle\psi|$ – *effect of state* $|\psi\rangle$ (dual vector, dual operator $V \rightarrow \mathbb{F}$ Hermitian conjugate) – bra
- *Inner product* of $|\psi\rangle$ and $|\varphi\rangle$ is $\langle\psi|\varphi\rangle$





Outer product

- *Outer product* $|w\rangle\langle v|$ for $w \in W, v \in V$ is a $V \rightarrow W$ operator:

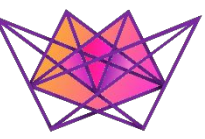
$$(|w\rangle\langle v|)(|v'\rangle) \equiv |w\rangle\langle v|v'\rangle = \langle v|v'\rangle|w\rangle$$

- Arbitrary $A: V \rightarrow W$ can be written

in a basis $\{|v_i\rangle\}_i$ for V and $\{|w_j\rangle\}_j$ for W :

$$A = \sum_{ij} m_{ij} |w_j\rangle\langle v_i|,$$

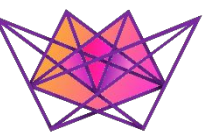
where $m_{ij} = \langle w_j|A|v_i\rangle$



Eigenvectors and eigenvalues

- $A = \sum_i \lambda_i |i\rangle\langle i|$, if $|i\rangle$ is an orthogonal basis in which A is diagonal.
- $|i\rangle$ – are *eigen vectors*
- λ_i – are *eigen values*
- Easy to check:

$$A|i\rangle = \lambda_i|i\rangle$$



Density operator (matrix)

- If $|\varphi_i\rangle$ – are pure states,
- $\{p_i\}$ – are probabilities over them, then

$$\rho \equiv \sum_i p_i |\varphi_i\rangle\langle\varphi_i| \text{ is a } \textit{dense operator}$$

- *Positive operator.*

$$\langle v|A|v\rangle \geq 0 \text{ for all } v$$

- **Theorem:** ρ is a density operator iff it's a positive Hermitian operator with trace = 1.



Trace inner product

- A and B are density matrices same dimension

- $tr(A^T B)$

- $A = \sum_i p_i |i\rangle\langle i|$ and $B = \sum_j q_j |j\rangle\langle j|$

- $tr(A^T B) = tr(\sum_i p_i |i\rangle\langle i| \sum_j q_j |j\rangle\langle j|) =$

$$= \sum_{ij} p_i q_j \langle i|j\rangle tr(|i\rangle\langle j|)$$

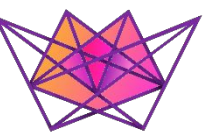
$$= \sum_{ij} p_i q_j \langle i|j\rangle \langle i|j\rangle$$

$$= \sum_{ij} p_i q_j \langle i|j\rangle^2$$



Distributional Semantics

- ‡ “You shall know the word by the company it keeps” (Firth)
 - Obtain meaning high dimensional vector representations from large corpora automatically
- Compositionality
 - DS can not be applied for entire sentence (lack of frequency)
- Entailment
 - w entails v if the meaning of a word w is included in the meaning of a word v
(w is-a v) – *subsumption relation*
 - non symmetric



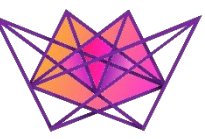
Distributional Inclusion Hypothesis

- If u is *semantically narrower* than v , then a significant number of salient distributional features of u are also included in the feature vector of v :
 - Hypothesis 1: If $v \Rightarrow w$ then all the characteristic features of v is expected to appear in w .
 - Hypothesis 2: If all the characteristic features of v appear in w , then $v \Rightarrow w$.



Category Theory

- A monoidal category \mathcal{C} is a category consisting of the following:
 - a functor $\otimes: \mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$ called the *tensor product*
 - an object $I \in \mathcal{C}$ called the *unit object*
 - a natural isomorphism whose components $(A \otimes B) \otimes C \xrightarrow{\alpha_{A,B,C}} A \otimes (B \otimes C)$ are called the *associators*
 - a natural isomorphism whose components $I \otimes A \xrightarrow{\lambda_A} A$ are called the *left unitors*
 - a natural isomorphism whose components $A \otimes I \xrightarrow{\rho_A} A$ are called the *right unitors*



Category Theory

- The objects of the category are thought to be types of systems
- A morphism $f: A \rightarrow B$ is a process that takes a system of type A to a system of type B .
- for $f: A \rightarrow B$ and $g: B \rightarrow C$, $g \circ f$ is the composite morphism that takes a system of type A into a system of type C by applying the process g after f .
- Morphisms of type $\psi: I \rightarrow A$ are called elements of A .



Compact closed categories

- A monoidal category is *compact closed* if for each object A , there are also left and right dual objects A^r and A^l , and morphisms

$$\eta^l: I \rightarrow A \otimes A^l \quad \eta^r: I \rightarrow A \otimes A^r$$

$$\epsilon^l: A^l \otimes A \rightarrow I \quad \epsilon^r: A \otimes A^r \rightarrow I$$

- that satisfies

$$(1_A \otimes \epsilon^l) \circ (\eta^l \otimes 1_A) = 1_A$$

$$(\epsilon^r \otimes 1_A) \circ (1_A \otimes \eta^r) = 1_A$$

$$(\epsilon^l \otimes 1_{A^l}) \circ (1_{A^l} \otimes \eta^l) = 1_{A^l}$$

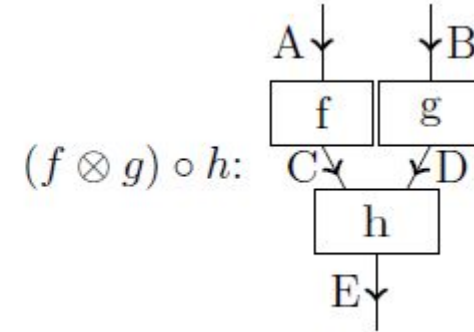
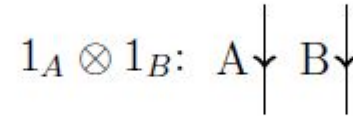
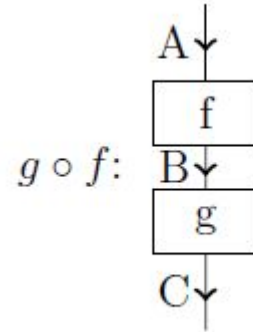
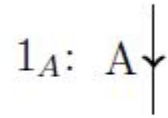
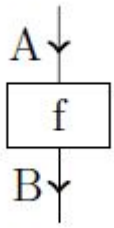
$$(1_{A^r} \otimes \epsilon^r) \circ (\eta^r \otimes 1_{A^r}) = 1_{A^r}$$

- The maps of compact categories are used to represent *correlations*, and in categorical quantum mechanics they model maximally entangled states.

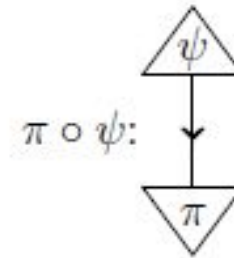
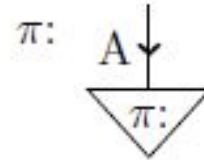
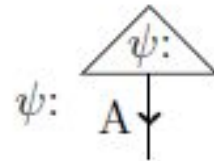


Graphical calculus

$f: A \rightarrow B$



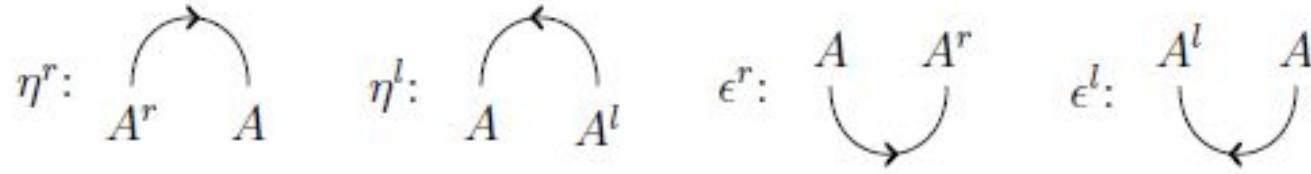
$\psi: I \rightarrow A$



$\pi: A \rightarrow I$



Graphical calculus



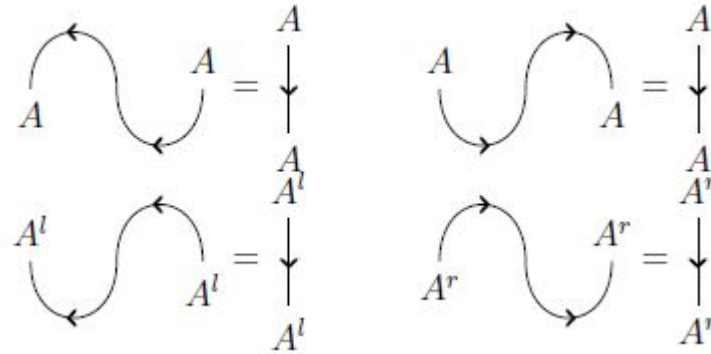
Snake identities

$$(1_A \otimes \epsilon^l) \circ (\eta^l \otimes 1_A) = 1_A$$

$$(\epsilon^r \otimes 1_A) \circ (1_A \otimes \eta^r) = 1_A$$

$$(\epsilon^l \otimes 1_{A^l}) \circ (1_{A^l} \otimes \eta^l) = 1_{A^l}$$

$$(1_{A^r} \otimes \epsilon^r) \circ (\eta^r \otimes 1_{A^r}) = 1_{A^r}$$



Swing rule





Compositional Distributional Model

- Pregroup grammars (Lambek)
- A partially ordered monoid $(P, \leq, \cdot, 1)$ consists of:
 - a set P
 - a monoid multiplication operator $\cdot : P \times P \rightarrow P$ satisfying the condition

$$(a \cdot b) \cdot c = a \cdot (b \cdot c) \text{ for all } a, b, c \in P$$

- and thy monoidal unit $1 \in P$ where for all $a \in P$
$$a \cdot 1 = a = 1 \cdot a$$
- a partial order \leq on P



Pregroup (Lambek, 2001)

- A *pregroup* $(P, \leq, \cdot, 1, (-)^l, (-)^r)$ is a partially ordered monoid in which each element a has both a left adjoint a^l and a right adjoint a^r such that

$$a^l a \leq 1 \leq a a^l \text{ and } a a^r \leq 1 \leq a^r a$$

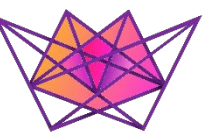
- Adjoints have properties:
 - Uniqueness: Adjoints are unique
 - Order reversal: If $a \leq b$ then $b^r \leq a^r$ and $b^l \leq a^l$
 - The unit is self adjoint: $1^l = 1 = 1^r$
 - Multiplication operation is self adjoint: $(a \cdot b)^l = b^l \cdot a^l$ and $(a \cdot b)^r = b^r \cdot a^r$
 - Opposite adjoints annihilate: $(a^r)^l = a = (a^l)^r$
 - Same adjoints iterate: $a^{ll} a^l \leq 1 \leq a^{rr} a^r, a^{lll} a^{ll} \leq 1 \leq a^{rrr} a^{rr}, \dots$



Pregroup grammar

- $a \rightarrow b$ means $a \leq b$ (a reduces to b)
- “John likes Mary”
- “John” and “Mary” assigned to type n (noun)
- “likes” is assigned to compound type $(n^r s n^l)$
- “likes” takes a noun from the left and from the right, and returns a sentence

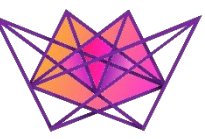
$$n(n^r s n^l)n \rightarrow 1s n^l n \rightarrow 1s1 \rightarrow s$$



Basic types

- n : noun
- j : infinitive of the verb

s : declarative statement
 g : glueing type



Pregroups as compact closed categories

- P is a concrete instance of a compact closed category

$$\eta^l = [1 \leq p \cdot p^l] \quad \epsilon^l = [p^l \cdot p \leq 1]$$

$$\eta^r = [1 \leq p^r \cdot p] \quad \epsilon^r = [p \cdot p^r \leq 1]$$

- Test snake identities:

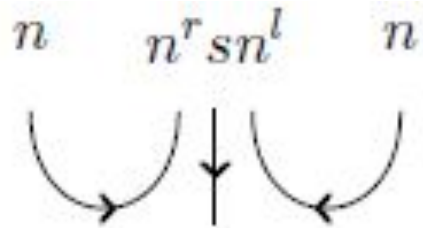
$$\begin{aligned} & (1^p \otimes \epsilon_p^l) \circ (\eta_p^l \otimes 1_p): \\ & p = 1p \leq pp^l p \leq p1 = p \end{aligned}$$

...



Examples

- “John likes Mary”



- “John does not like Mary”





FVect - finite dimensional vector space

- FVect – finite dimensional vector spaces over the base field \mathbb{R} together with linear maps, form a monoidal category
- FVect as a compact closed category.

Given a vector space V with basis $\{\vec{e}_i\}_i$

$$\eta_V^l = \eta_V^r : \mathbb{R} \rightarrow V \otimes V$$
$$1 \mapsto \sum_i e_i \otimes e_i$$

$$\epsilon_V^l = \epsilon_V^r : V \otimes V \rightarrow \mathbb{R}$$
$$\sum_{ij} c_{ij} v_i \otimes w_j \mapsto \sum_{ij} c_{ij} \langle v_i | w_j \rangle$$



$\mathbf{FVect} \times \mathbf{P}$ – categorical representation of meaning space

Objects in \mathbf{FVect} are of the form (V, p) , where V is the vector space representation of the meaning and p is the pregroup type. There exists a morphism $(f, \leq) : (V, p) \rightarrow (W, q)$ if there exists a morphism $f : V \rightarrow W$ in \mathbf{FVect} and $p \leq q$ in \mathbf{P} .

The compact closed structure of \mathbf{FVect} and \mathbf{P} lifts componentwise to the product category $\mathbf{FVect} \times \mathbf{P}$:

$$\eta^l : (\mathbb{R}, 1) \rightarrow (V \otimes V, p \cdot p^l)$$

$$\eta^r : (\mathbb{R}, 1) \rightarrow (V \otimes V, p^r \cdot p)$$

$$\epsilon^l : (V \otimes V, p^l \cdot p) \rightarrow (\mathbb{R}, 1)$$

$$\epsilon^r : (V \otimes V, p \cdot p^r) \rightarrow (\mathbb{R}, 1)$$

Definition 6.3. An object (V, p) in the product category is called a **meaning space**, where V is the vector space in which the meanings $\vec{v} \in V$ of strings of type p live.



“From-the-meanings-of-words-to-the-meanings-of-the-sentence” map

- Let $v_1 v_2 \dots v_n$ be a string of words, each v_i with a meaning space representation $\vec{v}_i \in (V_i, p_i)$. Let $x \in P$ be a pregroup type such that $[p_1 p_2 \dots p_n \leq x]$. Then the meaning vector for the string is:

$$\overrightarrow{v_1 v_2 \dots v_n} \in (W, x) \equiv f(v_1 \otimes v_2 \otimes \dots \otimes v_n),$$

- where f is defined to be the application of the compact closure maps obtained from the reduction $[p_1 p_2 \dots p_n \leq x]$ to the composite vector space $V_1 \otimes V_2 \otimes \dots \otimes V_n$.



Example: “John likes Mary”

- It has the pregroup type $nn^r sn^l n$
- vector representations $\overrightarrow{John}, \overrightarrow{Mary} \in V$ and $\overrightarrow{likes} \in V \otimes S \otimes V$
- The morphism in $\mathbf{FVect} \times \mathbf{P}$ corresponding to the map is of type:

$$(V \otimes (V \otimes S \otimes V) \otimes V, nn^r sn^l n) \rightarrow (s, S)$$

- From the pregroup reduction $[nn^r sn^l n \rightarrow s]$ we obtain the compact closure maps $\epsilon^r 1 \epsilon^l$. In \mathbf{FVect} this translates into:

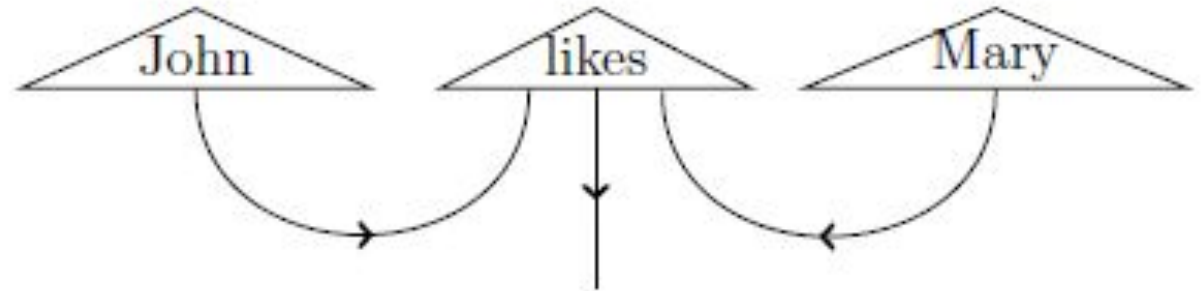
$$\epsilon_V \otimes 1_S \otimes \epsilon_V: V \otimes (V \otimes S \otimes V) \otimes V \rightarrow S$$



Example: “John likes Mary”

• $\overrightarrow{John} \otimes \overrightarrow{likes} \otimes \overrightarrow{Mary}$

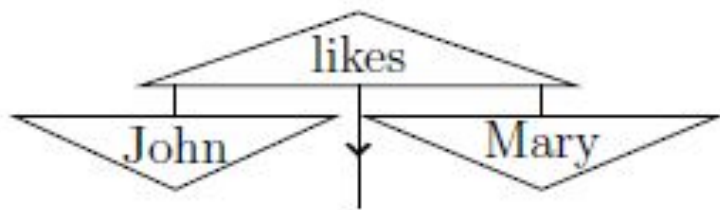
$$\overrightarrow{likes} = \sum_{ijk} c_{ijk} v_i \otimes s_j \otimes v_k$$



$$\overrightarrow{John\ likes\ Mary} = \epsilon_V \otimes 1_S \otimes \epsilon_V (\overrightarrow{John} \otimes \overrightarrow{likes} \otimes \overrightarrow{Mary})$$

$$= \sum_{ijk} \langle John | v_i \rangle s_j \langle v_k | Mary \rangle$$

$$(\langle \epsilon_V^r | \otimes 1_S \otimes \langle \epsilon_V^l |) \circ |\overrightarrow{John} \otimes \overrightarrow{likes} \otimes \overrightarrow{Mary}\rangle$$



$$(\langle \overrightarrow{John} | \otimes 1_S \otimes \langle \overrightarrow{Mary} |) \circ |\overrightarrow{likes}\rangle$$



Example

$$\begin{aligned} \lceil eat \rceil &= (|\overrightarrow{sloths}\rangle|\overrightarrow{plants}\rangle + |\overrightarrow{lions}\rangle|\overrightarrow{meat}\rangle)(\langle\overrightarrow{sloths}| \langle\overrightarrow{plants}| + \langle\overrightarrow{lions}| \langle\overrightarrow{meat}|) \\ &= (|\overrightarrow{sloths}\rangle|\overrightarrow{plants}\rangle)(\langle\overrightarrow{sloths}| \langle\overrightarrow{plants}|) + \\ &\quad (|\overrightarrow{sloths}\rangle|\overrightarrow{plants}\rangle)(\langle\overrightarrow{lions}| \langle\overrightarrow{meat}|) + \\ &\quad (|\overrightarrow{lions}\rangle|\overrightarrow{meat}\rangle)(\langle\overrightarrow{sloths}| \langle\overrightarrow{plants}|) + \\ &\quad (|\overrightarrow{lions}\rangle|\overrightarrow{meat}\rangle)(\langle\overrightarrow{lions}| \langle\overrightarrow{meat}|) \\ &\sim (|\overrightarrow{sloths}\rangle \langle\overrightarrow{sloths}| \otimes |\overrightarrow{plants}\rangle \langle\overrightarrow{plants}|) + \\ &\quad (|\overrightarrow{sloths}\rangle \langle\overrightarrow{lions}| \otimes |\overrightarrow{plants}\rangle \langle\overrightarrow{meat}|) + \\ &\quad (|\overrightarrow{lions}\rangle \langle\overrightarrow{sloths}| \otimes |\overrightarrow{meat}\rangle \langle\overrightarrow{plants}|) + \\ &\quad (|\overrightarrow{lions}\rangle \langle\overrightarrow{lions}| \otimes |\overrightarrow{meat}\rangle \langle\overrightarrow{meat}|) \end{aligned}$$



Example

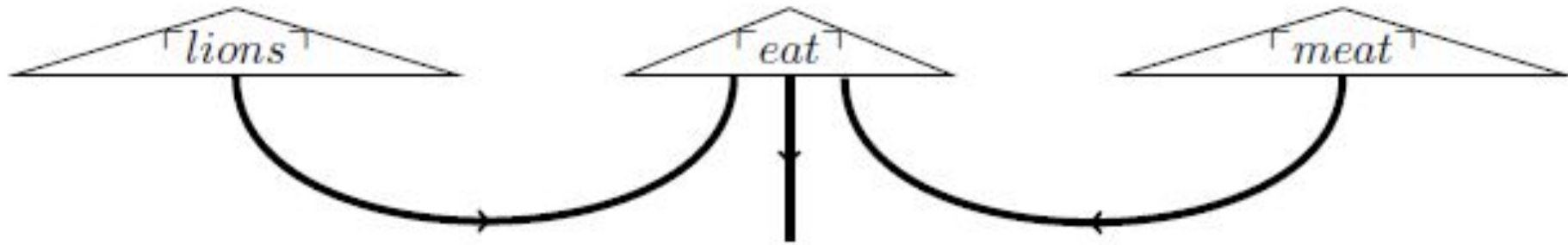
This is the density matrix representation of a pure composite state that relate “sloths” to “plants” and “lions” to “meat”. If we fix the bases $\{\overrightarrow{lions}, \overrightarrow{sloths}\}$ for N_1 , and $\{\overrightarrow{meat}, \overrightarrow{plants}\}$ for N_2 , $\lceil eat \rceil : N_1 \otimes N_1 \rightarrow N_2 \otimes N_2$ has the following matrix representation:

$$\begin{pmatrix} 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 \end{pmatrix}$$

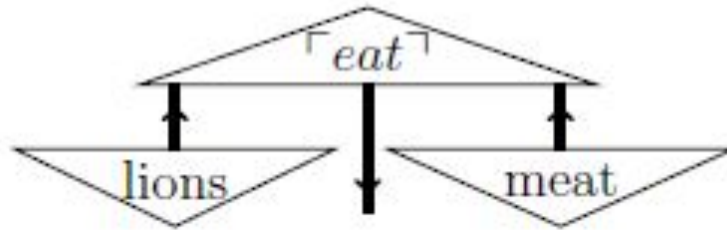
“Lions eat meat” . This is a transitive sentence, so as before, it gets assigned the pregroup type: $nn^l sn^r n$. The diagrammatic expression of the pregroup reduction is as follows:



Example



This reduces to:





Example

Explicit calculation gives:

$$\begin{aligned} & (\epsilon_N^l \otimes 1_S \otimes \epsilon_N^r)(\ulcorner lions \urcorner \otimes \ulcorner eat \urcorner \otimes \ulcorner meat \urcorner) \\ &= \langle \overrightarrow{lions} | \overrightarrow{sloths} \rangle^2 \langle \overrightarrow{plants} | \overrightarrow{meat} \rangle^2 + \\ & \quad \langle \overrightarrow{lions} | \overrightarrow{sloths} \rangle \langle \overrightarrow{lions} | \overrightarrow{lions} \rangle \langle \overrightarrow{meat} | \overrightarrow{meat} \rangle \langle \overrightarrow{plants} | \overrightarrow{meat} \rangle + \\ & \quad \langle \overrightarrow{lions} | \overrightarrow{lions} \rangle \langle \overrightarrow{lions} | \overrightarrow{sloths} \rangle \langle \overrightarrow{meat} | \overrightarrow{meat} \rangle \langle \overrightarrow{plants} | \overrightarrow{meat} \rangle + \\ & \quad \langle \overrightarrow{lions} | \overrightarrow{lions} \rangle^2 \langle \overrightarrow{meat} | \overrightarrow{meat} \rangle^2 \\ &= 0 + 0 + 0 + 1 \\ &= 1 \end{aligned}$$



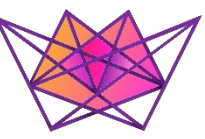
Example

“Sloths eat meat” . This sentence has a very similar calculation to the one above, and has the result:

$$(\epsilon_N^l \otimes 1_S \otimes \epsilon_N^r)(\ulcorner sloths \urcorner \otimes \ulcorner eat \urcorner \otimes \ulcorner meat \urcorner) = 0$$

“Mammals eat meat” . This sentence has the same pregroup types as the first sentence, and so has the same reduction map:

$$\begin{aligned} & (\epsilon_N^l \otimes 1_S \otimes \epsilon_N^r)(\ulcorner mammals \urcorner \otimes \ulcorner eat \urcorner \otimes \ulcorner meat \urcorner) \\ &= (\epsilon_N^l \otimes 1_S \otimes \epsilon_N^r)\left(\left(\frac{1}{2}\ulcorner lions \urcorner + \frac{1}{2}\ulcorner sloths \urcorner\right) \otimes \ulcorner eat \urcorner \otimes \ulcorner meat \urcorner\right) \\ &= \frac{1}{2}(\epsilon_N^l \otimes 1_S \otimes \epsilon_N^r)(\ulcorner lions \urcorner \otimes \ulcorner eat \urcorner \otimes \ulcorner meat \urcorner) + \\ & \quad \frac{1}{2}(\epsilon_N^l \otimes 1_S \otimes \epsilon_N^r)(\ulcorner sloths \urcorner \otimes \ulcorner eat \urcorner \otimes \ulcorner meat \urcorner) \\ &= \frac{1}{2} \end{aligned}$$



Readings

- Esma Balkir. Using Density Matrices in a Compositional Distributional Model of Meaning. // Master thesis. University of Oxford. 2014
- Joachim Lambek. Type grammars as pregroups. Grammars, 4(1):21{39, 2001.